Operating Systems

Tutorial 2 & 16

Michael Tänzer

os-tut@nhng.de http://os-tut.nhng.de

Calendar Week 44

Outline I

- Introduction
 - Organisational
 - Distro Wars
- Linux/*NIX Basics
 - The Shell
 - Man Pages
 - Find
 - Combining Commands
 - Build Process
- Bits & Bytes
- OS Basics
 - Policy vs. Mechanism
 - Major Tasks
 - Hardware Basics

Outline II

- User vs. Kernel Mode
- Memory Hierarchy
- Cache Organisation



About Me

Me, Myself and I

- Student of Computer Science (Diploma)
- 7th semester
- Interests:
 - Operating systems
 - Cryptography and security
 - Telematics

Introduction

0000000 Organisational

- Appear every two weeks
- 20 points each ⇒ 120 points total
- Have a deadline (don't miss it), usually Wednesday at 12:00 (noon)
- Will be marked and returned by me (not always on time)
- You may work in groups but everyone has to hand in separately
- If you only want to get your results and leave, you may do SO
- Otherwise please be patient until the end of the tutorial, when I will return them

Programming Assignments

Introduction

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#	Topic	Release	Deadline	Points
0	First Steps		None	0
1	Synchronisation	09.11.09	25.11.09	20
2	Memory Management	23.11.09	23.12.09	25
3	Syscalls	21.12.09	03.02.10	25

- You will work in pairs
- Register until 18.11.09 12:00 by sending an email to one of your tutors containing
 - Your names
 - Your matriculation numbers
 - A fancy nickname for your group
- Send your solution to the tutor chosen (compressed)
- We'll agree on a meeting where you show me what you've done

Programming Assignments

Theory

- Will not be marked
- Give an overview of the source needed for the assignment
- Basically you should do them but as you don't have to hand it in, you don't have to formulate the answer

Introduction

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- Do assignment 0 this will make sure you have a working system for the coming assignments
- Don't use qdb via sockets (as proposed in some eclipse) tutorials on the net) it's way too slow
- Use a VCS (Version Control System) to allow concurrent programming without headaches
 - I personally recommend git
 - I'll write a short git-Intro which you'll find on the website (http://os-tut.nhng.de - once it's done)

'Schein' & Bonus Points

Introduction

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- You can get up to 190 points
 - 120 from theory assignments
 - 70 from programming assignments
- For a 'Schein' you need ≥ 100 points
- If you are a bachelor student you need the 'Schein' to pass the module
- If you are a diploma student you don't need it
- Bonus points will be added to the points you get in your exam if you pass it \Rightarrow if you fail the exam the bonus points can't save you
 - \bullet > 110 points \Rightarrow 1 bonus point
 - \bullet > 130 points \Rightarrow 2 bonus point
 - \bullet > 150 points \Rightarrow 3 bonus point
 - \geq 170 points \Rightarrow 4 bonus point

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Now the fun part

Introduction

OOOOO

Distro Wars

- Your given name
- The Operating System and Distribution/Version you use most of the time
- For how long you've been using it

What is a Shell?

The shell is a CLI (Command Line Interface) for the Operating System

- In Linux the shell runs as a normal user programme
- Usually a command starts another programme which does the work
- The shell provides mechanisms to combine commands (pipes, control flow statements)
- And manipulate (environment) variables
- Allows manipulation of multiple files through wild cards
- Sophisticated shells (e. g. bash) also have convenience features like tab completion for almost everything

The Shell

Change the working directory cd

List the contents of a directory (like dir in Windows) 1.5

mkdir Make a new directory

Remove directory (only works if the directory is rmdir

empty)

Copy file or directory Ср

Move/Rename file or directory mν

Delete a file rm

Concatenate inputs to output cat.

less Scroll through input

Print the line containing a specified pattern in the grep

input

Execute a command with the input as parameters xarqs

Man Pages

How can you get help for a programme?

No, 'Use Google' is not the answer I've hoped for

- The command man shows the (hopefully complete and understandable) manual
- Sometimes more than one programme has the same name (e. g. there is the command, system call and function exit) then you also have to specify which section you want to see
- man man shows how to use man and which sections are available

Find

- Assume you want to edit a file named syscall.c, but you have forgotten in which subdirectory of our project it resides. What can you do?
- \$ find ./ -name 'syscall.c'
 ./kern/arch/mips/mips/syscall.c
 \$

Combine commands to get complex things done

- Assume you have a directory that contains multiple C source files in multiple subdirectories. How can you search for all occurrences of the macro FOOBAR in these C files?
- \$ find ./ -name '*.c'| xargs grep 'FOOBAR' ./foo/bar.c:#define FOOBAR 42
- \$ grep -r --include='*.c' 'FOOBAR'./
 ./foo/bar.c:#define FOOBAR 42

Creating Binaries

- What steps are necessary to create an executable program from multiple C source files?
 - Compilation that's gcc's job, each file is compiled on its own \Rightarrow multiple object files (*.o)
 - Object files contain machine code and not yet resolved symbols which reference to 'things' in other object files
 - Linkage that's what ld does, it resolves the references and puts everything together ⇒ one binary
- Usually gcc will also call 1d but you could give the -c option
- Manual linking is a little nightmare

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What is make?

- make is an automated build tool (an ancestor of ant)
- The build process is separated into steps called targets
- Targets can depend on other targets
- Reduces overhead:
 - Targets will only run once per invocation of make
 - Sources will only be compiled if they were modified since the last build
- Executes gcc and other tools to do the dirty work
- If something doesn't work try a make clean followed by make to get a clean build

Example Makefile

coffee: water powder filter

mv machine/water machine/filter/ && \

my machine/filter/water machine/coffee

water:

touch machine/water

powder: filter

touch machine/filter/powder

filter:

mkdir machine/filter

clean:

rm -r machine/*

Bit Operators

С	Description
&	Bitwise ∧
	Bitwise ∨
~	Bitwise ¬
<<	Right shift
>>	Left shift (filled with zeros)
^	Bitwise XOR

Why do we need bit operators?

- Why might it be necessary to set or clear a single bit of an integer value?
- Most hardware has control registers where each bit has a special meaning
- If one only needs to store one bit, why use more space?
 Bit maps, bit fields, etc.

Exercises

- How can you set the ith bit of a given integer value?
- result = value | (1 << i);
- How can you clear the ith bit of a given integer value?
- result = value & ~(1 << i);
- How can you retrieve the contents of bit 2 to bit 5 of a given integer value?
- result = value & (0xf << 2);

Policy vs. Mechanism

Difference between policy and mechanism

- Policies specify how a certain task should be fulfilled, optimised for certain goals.
- Mechanisms are used to realise those policies
- Example:
 - Task: Get from the Mensa to the 'Info-Bau'
 - Goals: Short, fast
 - Mechanisms: Turn left/right, slow down, etc.
 - Policies: Minimum deviation from linear distance, minimal number of obstacles

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Policy vs. Mechanism

Example II

- Assume you are the leader of a supermarket
- Task: How many cash points should be occupied?
- Goal: Minimal personnel costs, minimal waiting time for the customers
- Mechanisms:
 - Bell to indicate another cash point should be occupied
 - Icon on the cash display to indicate this cash point should be closed
 - Light barriers
 - Buzzers
- Policies: Have a minimum and maximum queue length, let the customers hit a buzzer if they wait too long

Enumerate the major tasks of an operating system

- Abstraction/Standardisation
- Resource Management
- Security/Protection
- Providing an execution environment for applications

User vs. Kernel Mode

Differences between a processor in user and kernel mode

Why are they needed?

- User mode:
 - Only unprivileged instructions may be executed
 - Otherwise Exception is thrown
- Kernel mode:
 - All instructions may be executed
 - Entered if an interrupt, trap or exception occurs
- Needed to protect applications from each other and make sure everyone gets his fair share

User vs. Kernel Mode

Typical examples for privileged instructions

Why are they privileged?

- Manipulate control registers for memory address translation
- Disable or enable interrupts
- Access privileged plattform devices

If they weren't user software could

- manipulate critical system state
- elevate its privileges
- overcome protection
- ruin isolation
- impact stability
- impact trustworthiness
- •

User vs. Kernel Mode

How could parameters for syscalls be passed to the kernel?

- Registers
 - Fast?
 - Limited number of registers
- Stack (or possibly another memory location)
 - Slower?
 - More flexible
 - Less hardware dependent

Principle and Benefits

- High speed memory is expensive
- Only the currently needed data is kept in high-speed memory
- If there is no space left for new data, data that hasn't been used for a long time is moved down in the hierarchy
- This is done recursively for L1, L2, (L3), RAM and hard disk
- Most of the time data can be found in high speed memory without a costly pure high-speed system

- Temporal locality: Data that has been accessed recently is likely to be accessed again (loops, etc.)
- Spatial locality: Data that is close together is likely to be accessed successively (arrays, etc.)
- Example: Stack

Memory Hierarchy

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Why is the cache divided into cache lines?

Which size?

- Transferring big parts at one time (burst mode) is faster than transferring multiple little units
- Transferring parts that are never used is unneeded effort
- ⇒ need to find the right balance

Cache Organisation

Fully Associative

Memory Address Tag **Block**

- Each memory block may be mapped to each cache line
- Ideal utilisation of the given space
- Lots of comparators needed (one per cache line) ⇒ expensive

Block

Cache Organisation

Direct Mapped

Memory Address Tag Cache Line

- Each memory block is mapped on the cache line corresponding to the lower bits of its address (without the bits used within each block)
- Only one comparator needed
- If more than one memory block that are mapped to the same cache line are in use they replace each other ⇒ cache useless

n-way Set Associative

Cache Organisation

Memory Address Tag | Set | Block

- Compromise of the other two
- Each memory block is mapped to a subset of n cache lines (like the direct mapped approach) inside these sets it's mapped like in the fully associative organisation
- + n comparators needed
- + It takes more than *n* memory blocks mapped to the same set to cause the blocks replacing each other
- $n = 1 \Rightarrow$ direct mapped, $n = m \Rightarrow$ fully associative

Questions & Comments

Any questions or comments?

A Quick Survey

Write on an anonymous piece of paper:

- At least one thing you liked
- At least one thing that could be improved about the tutorial

The End

The End